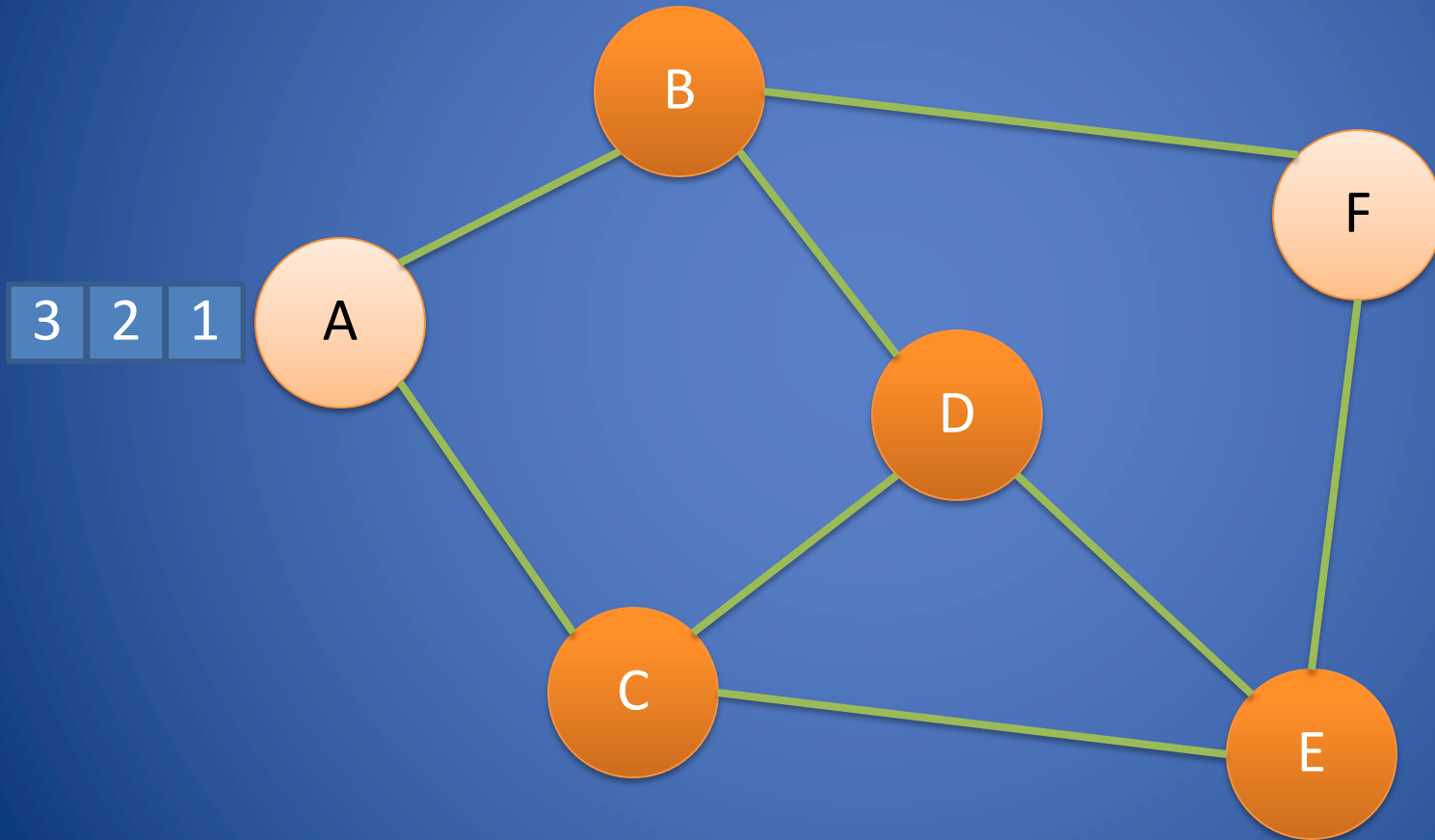


Computer Networks

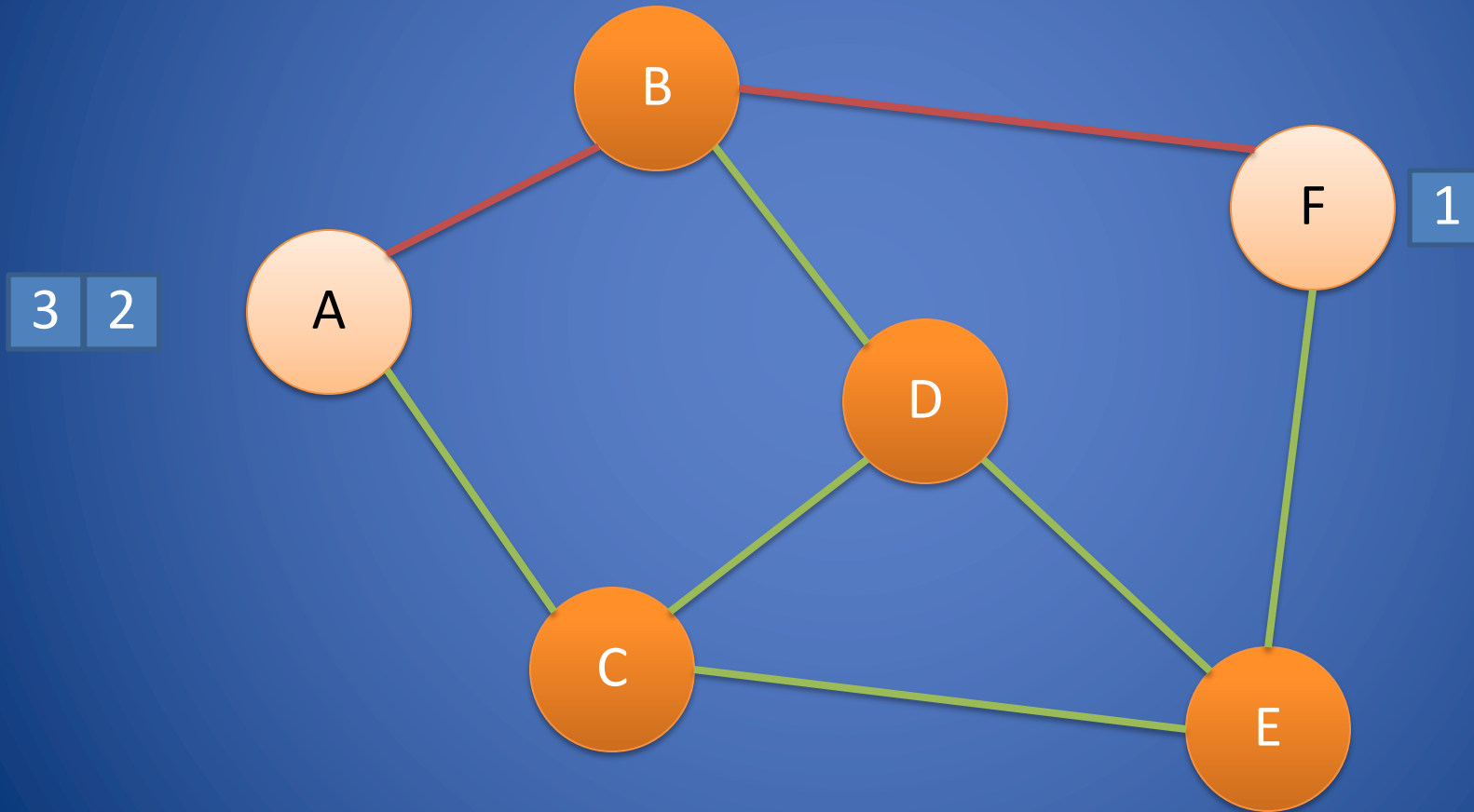
Circuit and Packet Switching

- Circuit switching
 - Legacy phone network
 - Single route through sequence of hardware devices established when two nodes start communication
 - Data sent along route
 - Route maintained until communication ends
- Packet switching
 - Internet
 - Data split into **packets**
 - Packets transported independently through network
 - Each packet handled on a **best efforts** basis
 - Packets may follow different routes

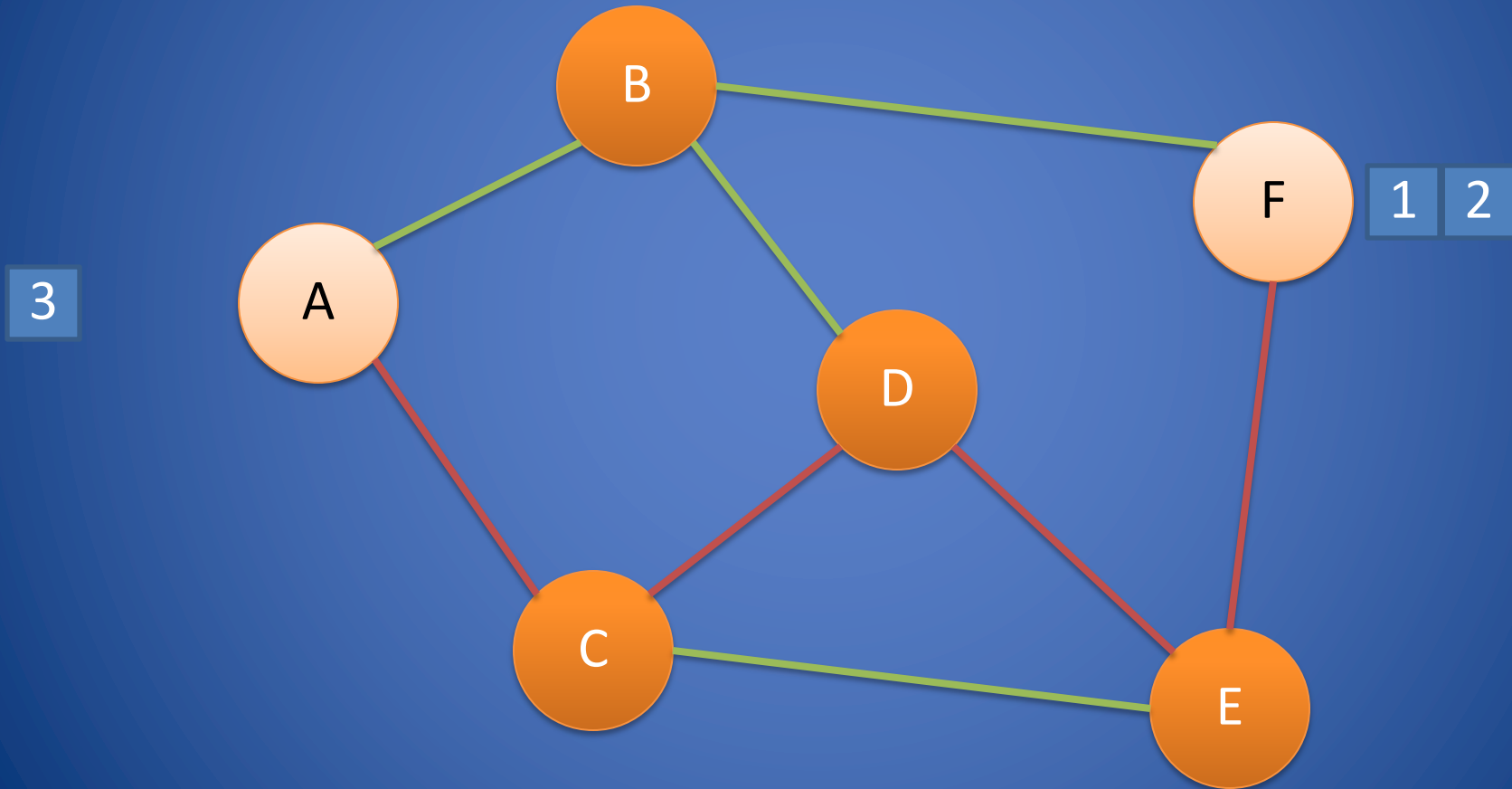
Packet Switching



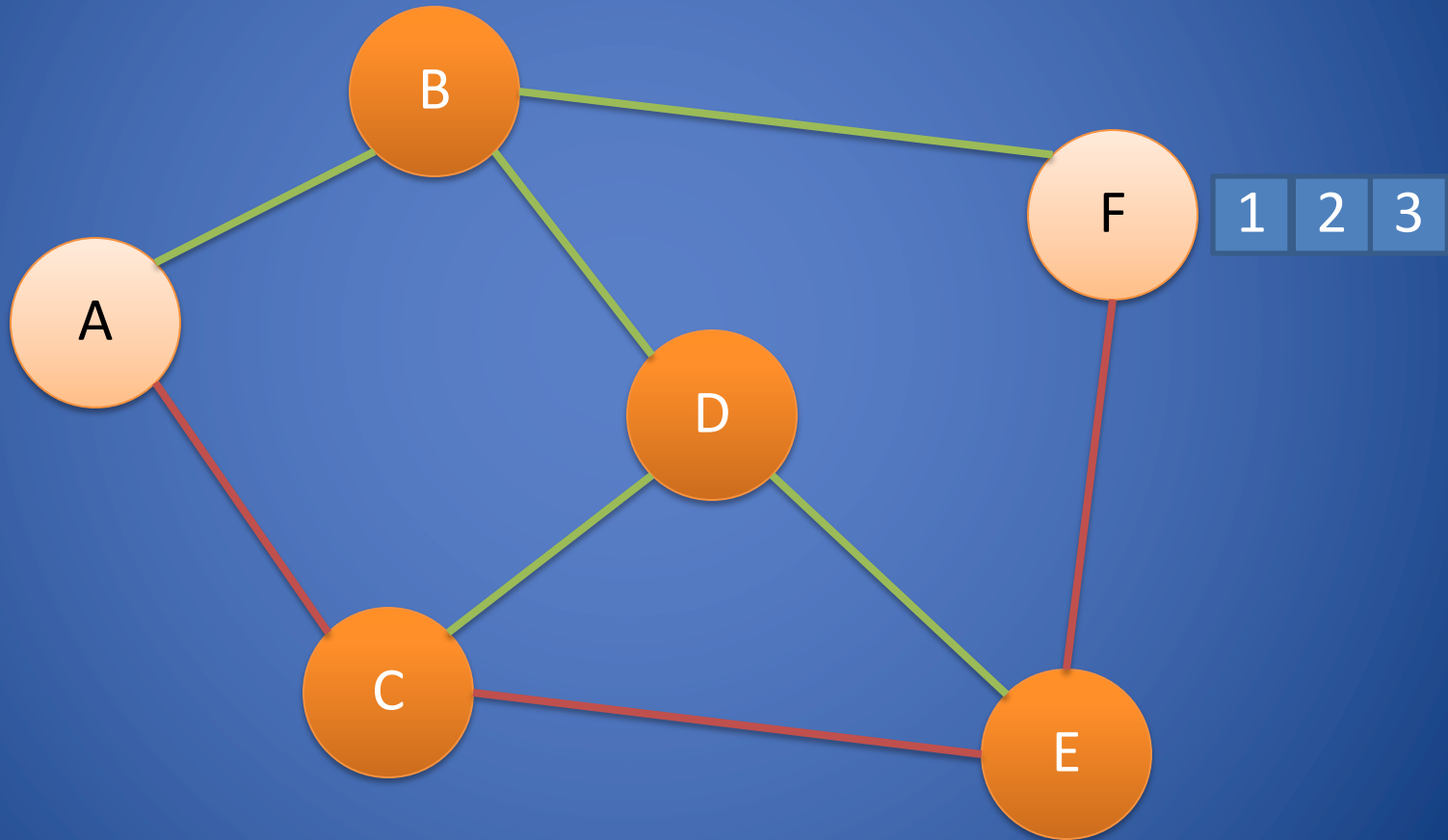
Packet Switching



Packet Switching



Packet Switching



Protocols

- A **protocol** defines the rules for communication between computers
- Protocols are broadly classified as connectionless and connection oriented
- **Connectionless protocol**
 - Sends data out as soon as there is enough data to be transmitted
 - E.g., user datagram protocol (UDP)
- **Connection-oriented protocol**
 - Provides a reliable connection stream between two nodes
 - Consists of set up, transmission, and tear down phases
 - Creates virtual circuit-switched network
 - E.g., transmission control protocol (TCP)

Encapsulation

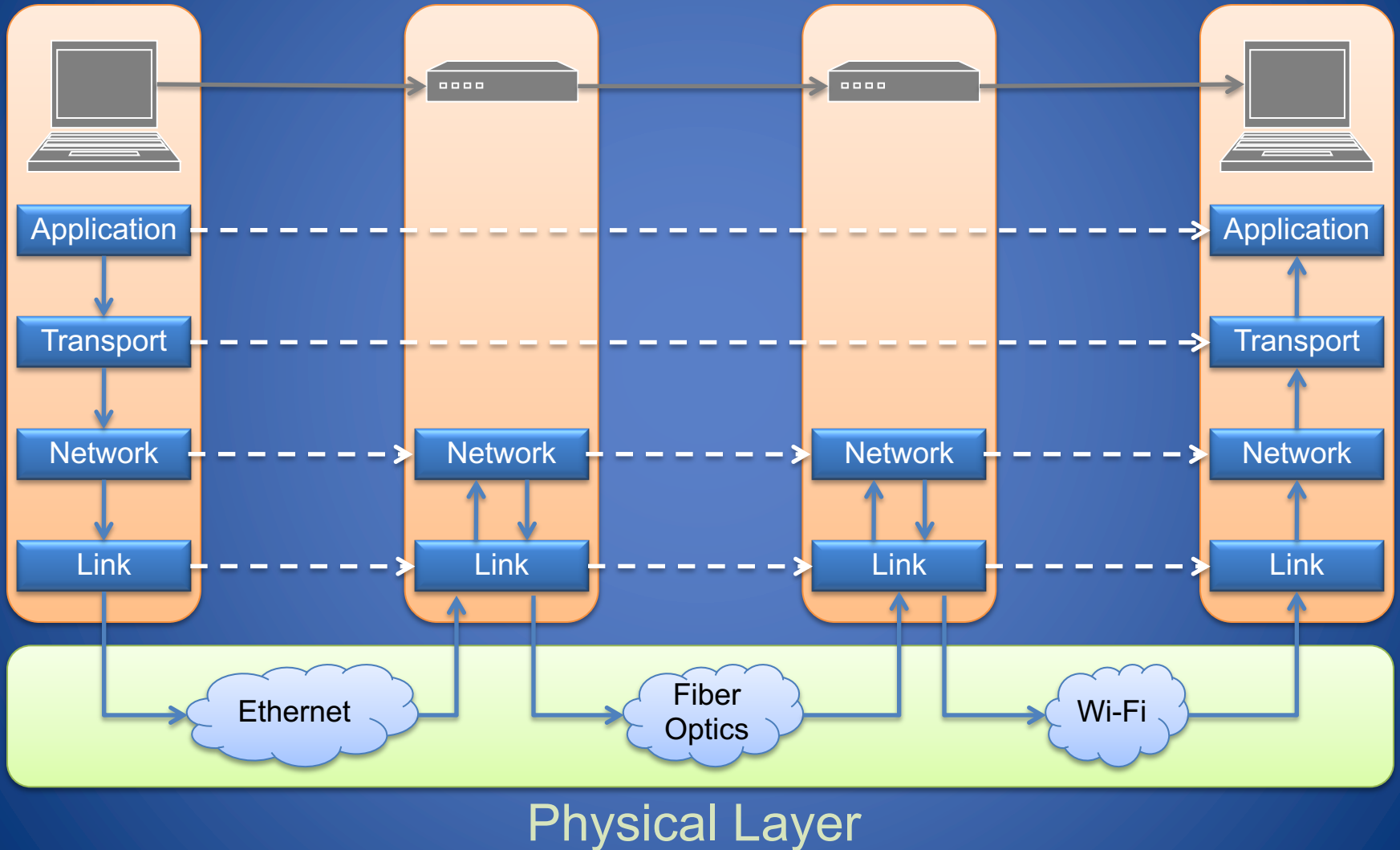
- A packet typically consists of
 - Control information for addressing the packet: **header** and **footer**
 - Data: **payload**
- A network protocol N1 can use the services of another network protocol N2
 - A packet p1 of N1 is encapsulated into a packet p2 of N2
 - The payload of p2 is p1
 - The control information of p2 is derived from that of p1



Network Layers

- Network models typically use a **stack** of layers
 - Higher layers use the services of lower layers via encapsulation
 - A layer can be implemented in hardware or software
 - The bottommost layer must be in hardware
- A network device may implement several layers
- A communication channel between two nodes is established for each layer
 - Actual channel at the bottom layer
 - Virtual channel at higher layers

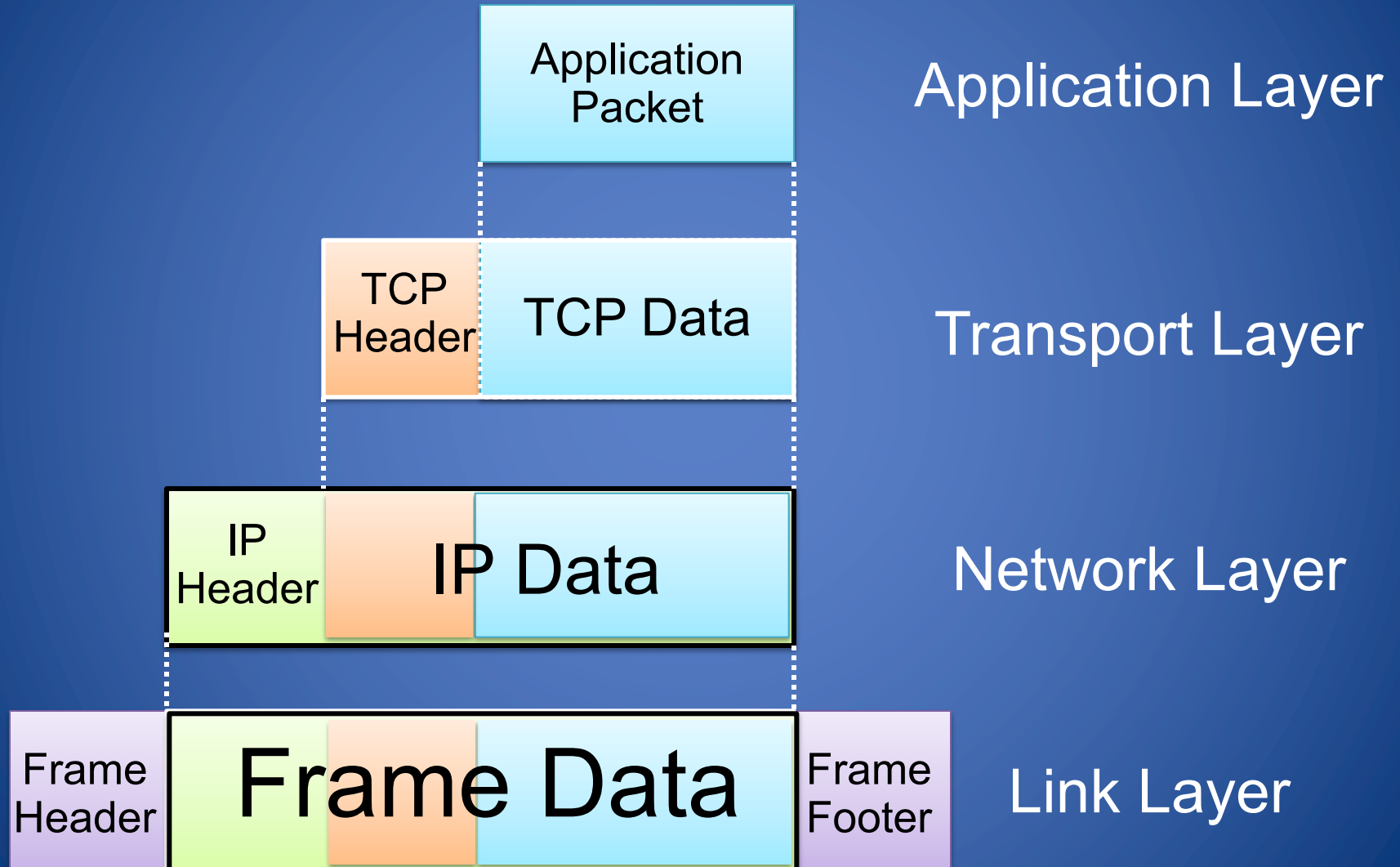
Internet Layers



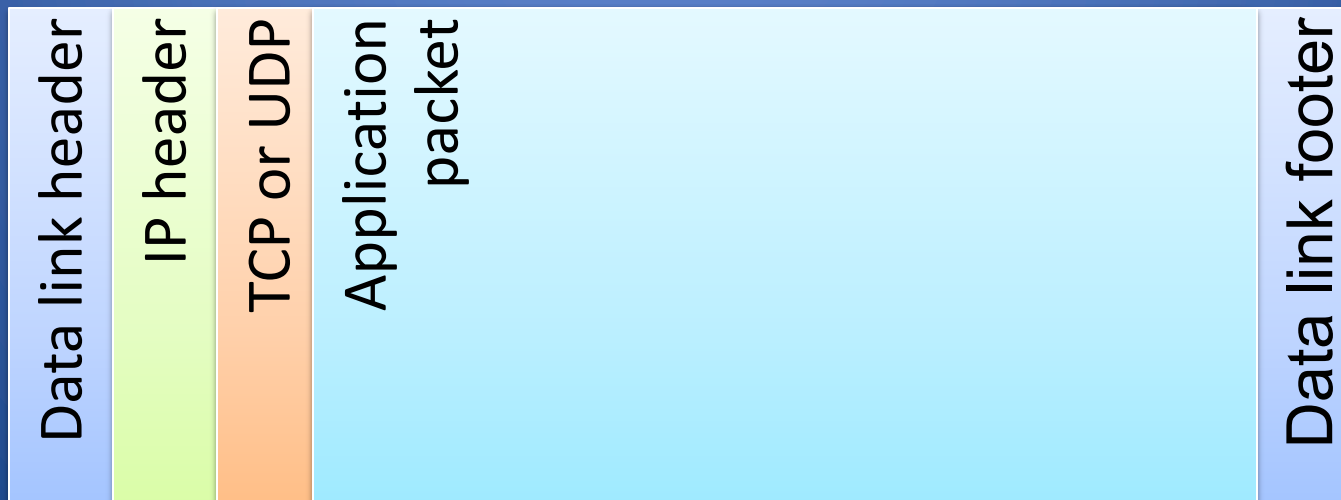
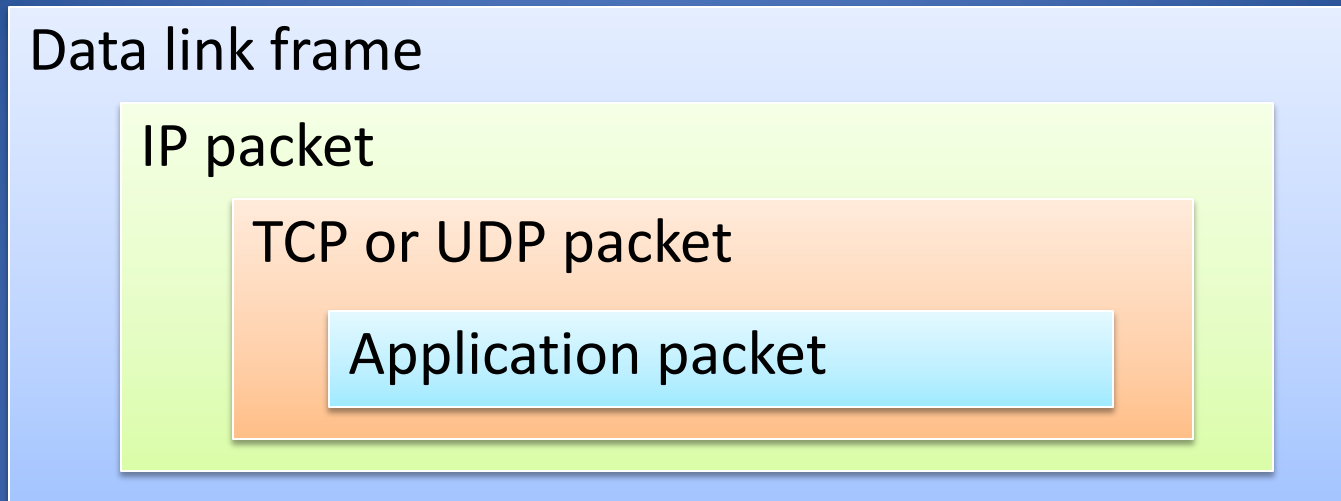
Intermediate Layers

- Link layer
 - Local area network: Ethernet, WiFi, optical fiber
 - 48-bit media access control (**MAC**) addresses
 - Packets called **frames**
- Network layer
 - Internet-wide communication
 - Best efforts
 - 32-bit internet protocol (**IP**) addresses in IPv4
 - 128-bit IP addresses in IPv6
- Transport layer
 - 16-bit addresses (**ports**) for classes of applications
 - Connection-oriented transmission layer protocol (**TCP**)
 - Connectionless user datagram protocol (**UDP**)

Internet Packet Encapsulation

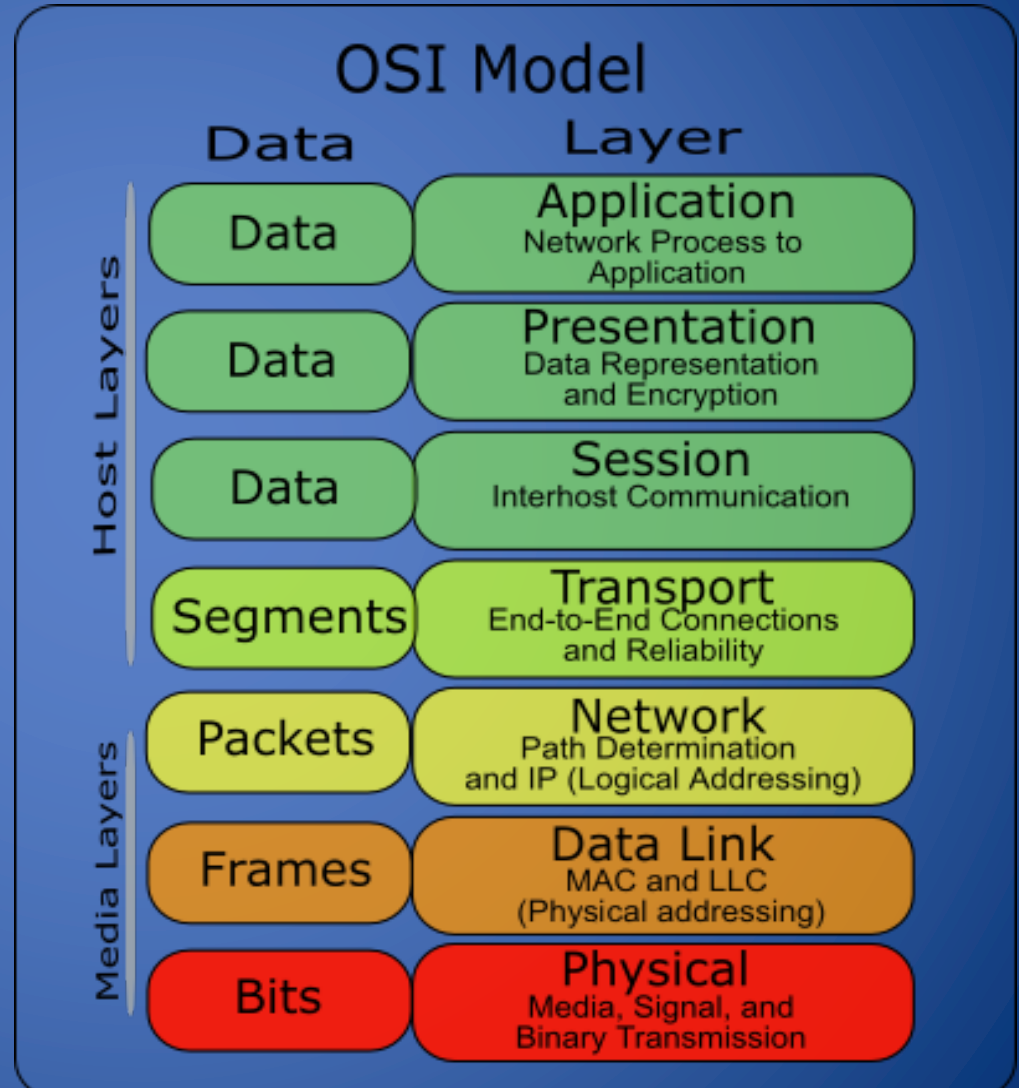


Internet Packet Encapsulation



The OSI Model

- The OSI (Open System Interconnect) Reference Model is a network model consisting of seven layers
- Created in 1983, OSI is promoted by the International Standard Organization (ISO)



Network Interfaces

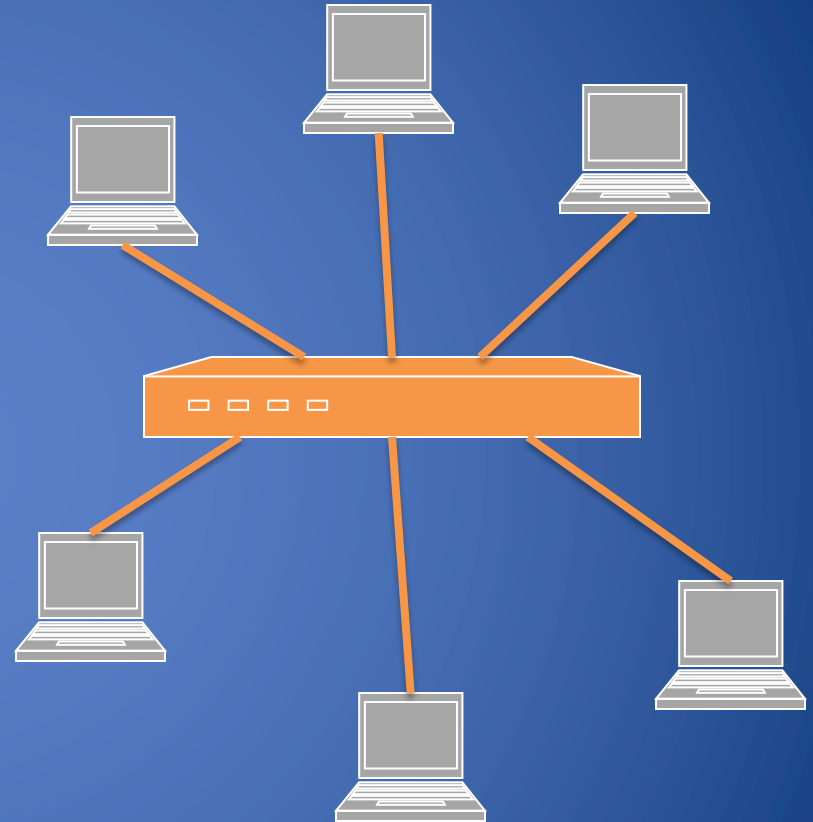
- Network interface: device connecting a computer to a network
 - Ethernet card
 - WiFi adapter
- A computer may have multiple network interfaces
- Packets transmitted between network interfaces
- Most local area networks, (including Ethernet and WiFi) broadcast frames
- In regular mode, each network interface gets the frames intended for it
- Traffic sniffing can be accomplished by configuring the network interface to read all frames (**promiscuous mode**)

MAC Addresses

- Most network interfaces come with a predefined MAC address
- A MAC address is a 48-bit number usually represented in hex
 - E.g., 00-1A-92-D4-BF-86
- The first three octets of any MAC address are IEEE-assigned Organizationally Unique Identifiers
 - E.g., Cisco 00-1A-A1, D-Link 00-1B-11, ASUSTek 00-1A-92
- The next three can be assigned by organizations as they please, with uniqueness being the only constraint
- Organizations can utilize MAC addresses to identify computers on their network
- MAC address can be reconfigured by network interface driver software

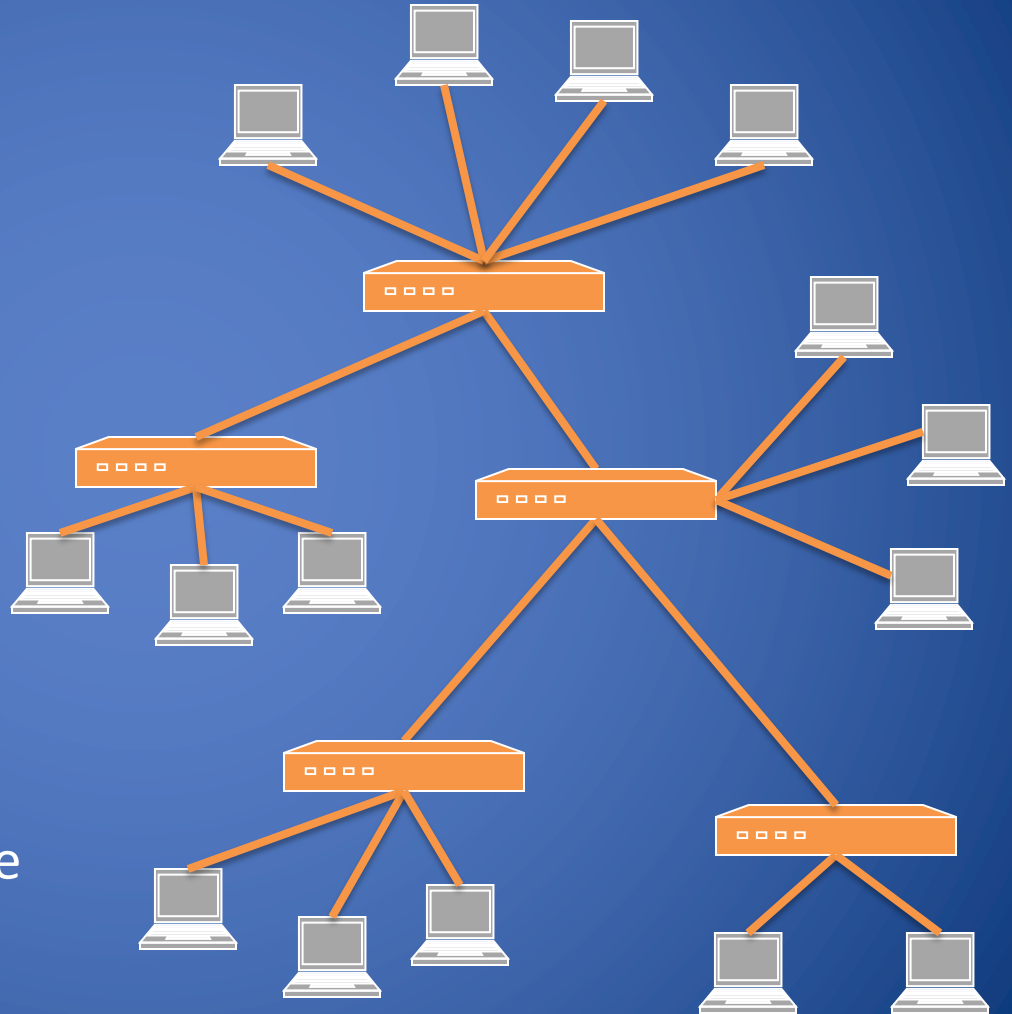
Switch

- A **switch** is a common network device
 - Operates at the link layer
 - Has multiple ports, each connected to a computer
- Operation of a switch
 - Learn the MAC address of each computer connected to it
 - Forward frames only to the destination computer



Combining Switches

- Switches can be arranged into a **tree**
- Each port learns the MAC addresses of the machines in the segment (subtree) connected to it
- Fragments to unknown MAC addresses are broadcast
- Frames to MAC addresses in the same segment as the sender are ignored



MAC Address Filtering

- A switch can be configured to provide service only to machines with specific MAC addresses
- Allowed MAC addresses need to be registered with a network administrator
- A MAC spoofing attack impersonates another machine
 - Find out MAC address of target machine
 - Reconfigure MAC address of rogue machine
 - Turn off or unplug target machine
- Countermeasures
 - Block port of switch when machine is turned off or unplugged
 - Disable duplicate MAC addresses

Viewing and Changing MAC Addresses

- Viewing the MAC addresses of the interfaces of a machine
 - Linux: `ifconfig`
 - Windows: `ipconfig /all`
- Changing a MAC address in Linux
 - Stop the networking service: `/etc/init.d/network stop`
 - Change the MAC address: `ifconfig eth0 hw ether <MAC-address>`
 - Start the networking service: `/etc/init.d/network start`
- Changing a MAC address in Windows
 - Open the Network Connections applet
 - Access the properties for the network interface
 - Click “Configure ...”
 - In the advanced tab, change the network address to the desired value
- Changing a MAC address requires administrator privileges

ARP

- The **address resolution protocol (ARP)** connects the network layer to the data layer by converting IP addresses to MAC addresses
- ARP works by **broadcasting** requests and caching responses for future use
- The protocol begins with a computer broadcasting a message of the form
who has <IP address1> tell <IP address2>
- When the machine with **<IP address1>** or an ARP server receives this message, it broadcasts the response
<IP address1> is <MAC address>
- The requestor's IP address **<IP address2>** is contained in the link header
- The Linux and Windows command **arp - a** displays the ARP table

Internet Address	Physical Address	Type
128.148.31.1	00-00-0c-07-ac-00	dynamic
128.148.31.15	00-0c-76-b2-d7-1d	dynamic
128.148.31.71	00-0c-76-b2-d0-d2	dynamic
128.148.31.75	00-0c-76-b2-d7-1d	dynamic
128.148.31.102	00-22-0c-a3-e4-00	dynamic
128.148.31.137	00-1d-92-b6-f1-a9	dynamic

ARP Spoofing

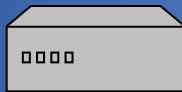
- The ARP table is updated whenever an ARP response is received
- Requests are not tracked
- ARP announcements are not authenticated
- Machines trust each other
- A rogue machine can spoof other machines

ARP Poisoning (ARP Spoofing)

- According to the standard, almost all ARP implementations are stateless
- An arp cache updates every time that it receives an arp reply... even if it did not send any arp request!
- It is possible to “poison” an arp cache by sending **gratuitous arp replies**
- Using static entries solves the problem but it is almost impossible to manage!

ARP Caches

IP: 192.168.1.1
MAC: 00:11:22:33:44:01



Data

IP: 192.168.1.105
MAC: 00:11:22:33:44:02



192.168.1.1 is at
00:11:22:33:44:01

192.168.1.105 is at
00:11:22:33:44:02

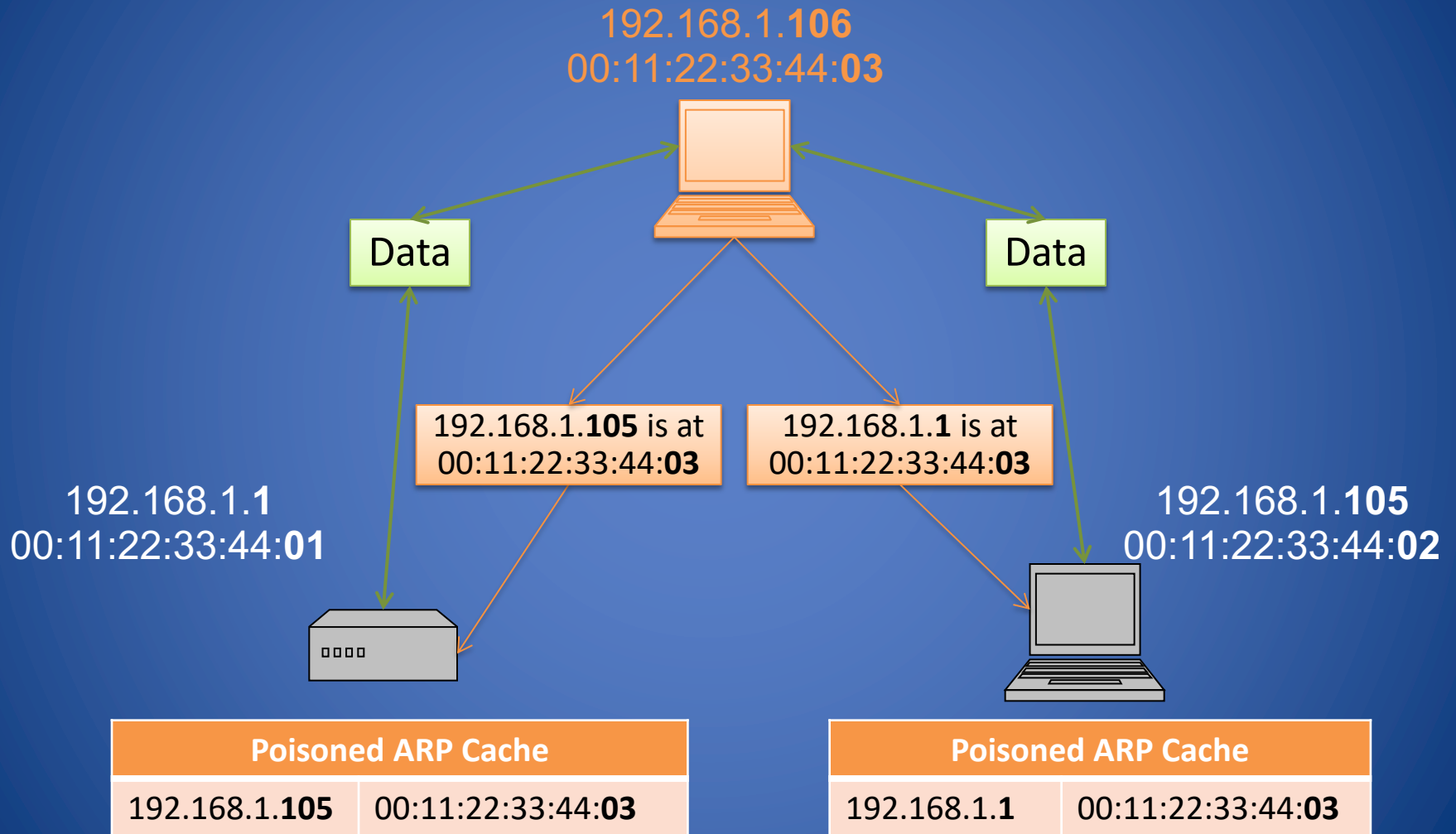
ARP Cache

192.168.1.105	00:11:22:33:44:02
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ARP Cache

192.168.1.1	00:11:22:33:44:01
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Poisoned ARP Caches



Root cause and defense

- The ARP spoofing is derived from the lack of identity verification in the Internet's underlying mechanisms.
- Defense:
 - Checking for multiple occurrences of the same MAC address on the LAN.
 - Manually specify a router's ARP cache to assign certain MAC addresses to specify IP addresses. Requires to adjust the cache are ignored.